

## Vježbe 10

1. (a) Odrediti generatorski polinom za BCH kodove (15,2) i (15,3), koristeći prosti polinom  $x^4 + x^3 + 1$ .
- (b) Data je poruka 10111. Kodirati je BCH kodom (15,3), definisanim pod (a).

### Rješenje:

- (a) Označimo sa  $\alpha$  korjenu posmatranog prostog polinoma  $p(x) = x^4 + x^3 + 1$ . Dalje možemo pisati:

$$\alpha^4 + \alpha^3 + 1 = 0, \text{ odnosno } \alpha^4 = \alpha^3 + 1.$$

Prikažimo sada svaki stepen  $\alpha^i$  pomoću izvedene jednakosti:

$$\alpha^5 = \alpha\alpha^4 = \alpha(\alpha^3 + 1) = \alpha^4 + \alpha = \alpha^3 + \alpha + 1,$$

$$\alpha^6 = \alpha\alpha^5 = \alpha(\alpha^3 + \alpha + 1) = \alpha^4 + \alpha^2 + \alpha = \alpha^3 + \alpha^2 + \alpha + 1,$$

$$\alpha^7 = \alpha\alpha^6 = \alpha(\alpha^3 + \alpha^2 + \alpha + 1) = \alpha^4 + \alpha^3 + \alpha^2 + \alpha = \alpha^2 + \alpha + 1,$$

$$\alpha^8 = \alpha\alpha^7 = \alpha(\alpha^2 + \alpha + 1) = \alpha^3 + \alpha^2 + \alpha,$$

$$\alpha^9 = \alpha\alpha^8 = \alpha(\alpha^3 + \alpha^2 + \alpha) = \alpha^4 + \alpha^3 + \alpha^2 = \alpha^2 + 1,$$

$$\alpha^{10} = \alpha\alpha^9 = \alpha(\alpha^2 + 1) = \alpha^3 + \alpha,$$

$$\alpha^{11} = \alpha\alpha^{10} = \alpha(\alpha^3 + \alpha) = \alpha^4 + \alpha^2 = \alpha^3 + \alpha^2 + 1,$$

$$\alpha^{12} = \alpha\alpha^{11} = \alpha(\alpha^3 + \alpha^2 + 1) = \alpha^4 + \alpha^3 + \alpha = \alpha + 1,$$

$$\alpha^{13} = \alpha\alpha^{12} = \alpha(\alpha + 1) = \alpha^2 + \alpha,$$

$$\alpha^{14} = \alpha\alpha^{13} = \alpha(\alpha^2 + \alpha) = \alpha^3 + \alpha^2,$$

$$\alpha^{15} = \alpha\alpha^{14} = \alpha(\alpha^3 + \alpha^2) = \alpha^4 + \alpha^3 = \alpha^3 + 1 + \alpha^3 = 1.$$

Predstavimo dobijene rezultate tabelarno, radi preglednosti:

$\alpha^0$	0001
$\alpha^1$	0010
$\alpha^2$	0100
$\alpha^3$	1000
$\alpha^4$	1001
$\alpha^5$	1011
$\alpha^6$	1111
$\alpha^7$	0111
$\alpha^8$	1110
$\alpha^9$	0101
$\alpha^{10}$	1010
$\alpha^{11}$	1101
$\alpha^{12}$	0011
$\alpha^{13}$	0110
$\alpha^{14}$	1100

Konjugati su stepeni reda  $2^n$  stepena korjena prostog polinoma, koji se koriste u konstrukciji koda. Pošto konstruišemo kodove koji maksimalno mogu ispraviti tri greške,  $w=3$ , u našem slučaju to su stepeni:

$$\alpha^{2^{w-1}} \text{ za } w=1,2,3, \text{ odnosno } [\alpha, \alpha^3, \alpha^5].$$

Odgovarajući konjugati za ove stepene su:

$$\alpha : \alpha, \alpha^2, \alpha^4, \alpha^8;$$

$$\alpha^3 : \alpha^3, \alpha^6, \alpha^9, \alpha^{12};$$

$$\alpha^5 : \alpha^5, \alpha^{10}.$$

Obratite pažnju na to da se konjugati se ne ponavljaju, odnosno da se uzima samo broj različitih vrijednosti konjugata!

Npr. kada tražimo konjugate za stepen  $\alpha$ :

$$\alpha^{2^n} \text{ za } n=0 \text{ dobijamo } \alpha,$$

$$\text{za } n=1 \text{ dobijamo } \alpha^2,$$

$$\text{za } n=2 \text{ dobijamo } \alpha^4,$$

$$\text{za } n=3 \text{ dobijamo } \alpha^8,$$

$$\text{za } n=4 \text{ dobijamo } \alpha^{16} = \alpha\alpha^{15} = \alpha \text{ (sjetite se da smo dokazali da je } \alpha^{15} = 1 \text{ !), a } \alpha \text{ već imamo!}$$

Pošto smo sa  $\alpha$  označili korjene posmatranog prostog polinoma, definišimo i preostala dva polinoma (za stepene  $\alpha^3$  i  $\alpha^5$ ). Označimo ih sa  $g_3(x)$  i  $g_5(x)$ .

Pri tome važi  $g_3(\alpha^3)=0$  i  $g_5(\alpha^5)=0$  (isto tako važi za naš prosti polinom  $p(\alpha)=0$ ).

$$g_3(x) = g_{30} + g_{31}x + g_{32}x^2 + g_{33}x^3 + g_{34}x^4$$

$$g_{30}[0001] + g_{31}[1000] + g_{32}[1111] + g_{33}[0101] + g_{34}[0011] = [0000]$$

$$\alpha^3 \quad \alpha^6 \quad \alpha^9 \quad \alpha^{12}$$

Napišimo sada odgovarajući sistem jednačina, u cilju dobijanja koeficijenata ovog polinoma:

$$g_{31} + g_{32} = 0$$

$$g_{32} + g_{33} = 0$$

$$g_{32} + g_{34} = 0$$

$$g_{30} + g_{32} + g_{33} + g_{34} = 0$$

Netrivijalno rješenje je:  $g_{30} = g_{31} = g_{32} = g_{33} = g_{34} = 1$ , pa je:

$$g_3(x) = x^4 + x^3 + x^2 + x + 1.$$

Na sličan način odredimo i polinom  $g_5(x)$ :

$$g_5(x) = g_{50} + g_{51}x + g_{52}x^2$$

$$g_{50}[0001] + g_{51}[1011] + g_{52}[1010] = [0000]$$

$$\alpha^5 \quad \alpha^{10}$$

Odgovarajući sistem jednačina je:

$$g_{51} + g_{52} = 0$$

$$g_{51} + g_{52} = 0$$

$$g_{50} + g_{51} = 0$$

Netrivijalno rješenje je:  $g_{50} = g_{51} = g_{52} = 1$ , pa je:

$$g_5(x) = x^2 + x + 1.$$

Napišimo sada generatorski polinom za BCH kod (15,2):

$$g(x) = p(x)g_3(x) = (x^4 + x^3 + 1)(x^4 + x^3 + x^2 + x + 1) = x^8 + x^4 + x^2 + x + 1.$$

Slično, generatorski polinom za BCH kod (15,3):

$$g(x) = p(x)g_3(x)g_5(x) = (x^4 + x^3 + 1)(x^4 + x^3 + x^2 + x + 1)(x^2 + x + 1) = x^{10} + x^9 + x^8 + x^6 + x^5 + x^2 + 1.$$

(b) Kodiranu poruku dobijamo množenjem informacionih simbola i generatorskog polinoma:

$$c(x) = i(x)g(x).$$

Riječ koju treba da kodiramo je 10111, pa je njen odgovarajući polinom u opštem obliku:

$$i(x) = i_0 + i_1x + i_2x^2 + i_3x^3 + i_4x^4.$$

Odredimo sada odgovarajuću kodiranu riječ:

$$\begin{aligned} c(x) &= (i_0 + i_1x + i_2x^2 + i_3x^3 + i_4x^4)(x^{10} + x^9 + x^8 + x^6 + x^5 + x^2 + 1) = \\ &= i_0 + i_1x + (i_0 + i_2)x^2 + (i_1 + i_3)x^3 + (i_2 + i_4)x^4 + \\ &+ (i_0 + i_3)x^5 + (i_0 + i_1 + i_4)x^6 + (i_1 + i_2)x^7 + (i_0 + i_2 + i_3)x^8 + (i_0 + i_1 + i_3 + i_4)x^9 + \\ &+ (i_0 + i_1 + i_2 + i_4)x^{10} + (i_1 + i_2 + i_4)x^{11} + (i_2 + i_3 + i_4)x^{12} + (i_3 + i_4)x^{13} + i_4x^{14} \end{aligned}$$

$$\begin{array}{ll} c_0 = i_0 & c_8 = i_0 + i_2 + i_3 \\ c_1 = i_1 & c_9 = i_0 + i_1 + i_3 + i_4 \\ c_2 = i_0 + i_2 & c_{10} = i_0 + i_1 + i_2 + i_4 \\ c_4 = i_1 + i_3 & c_{11} = i_1 + i_2 + i_4 \\ c_5 = i_0 + i_3 & c_{12} = i_2 + i_3 + i_4 \\ c_6 = i_0 + i_1 + i_4 & c_{13} = i_3 + i_4 \\ c_7 = i_1 + i_2 & c_{14} = i_4 \end{array}$$

Kodirana riječ je:

$$\begin{array}{cccccccccccccccc} c(x) = & 1 & 0 & 0 & 1 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 0 & 1 & 0 & 1 \\ & c_0 & c_1 & c_2 & c_3 & c_4 & c_5 & c_6 & c_7 & c_8 & c_9 & c_{10} & c_{11} & c_{12} & c_{13} & c_{14} \end{array}$$